

**Coffman, E. G. jun.; Gilbert, E. N.**

**On the expected relative performance of list scheduling.** (English) Zbl 0569.90044  
Oper. Res. 33, 548-561 (1985).

Let  $\bar{X} = (X_1, \dots, X_n)$  denote an ordered list of service times required by  $n$  tasks. The service will be performed by  $m \geq 2$  processors working in parallel. Each processor serves one task at a time and, having once started a task, finishes it before starting another. A schedule determines how the tasks are to be served. A list schedule keeps the tasks not yet serviced listed in the order prescribed by  $\bar{X}$ . Whenever a processor completes a service, it then takes its next task from the head of the list. The makespan of a schedule is the time required for all service to be completed. The makespan  $L(\bar{X})$  of a list schedule is usually longer than necessary. Reordering the tasks in an optimal way can reduce the makespan to  $OPT(\bar{X})$ , the smallest possible makespan, but requires knowing the  $X_i$  in advance and solving an NP-complete problem. The ratio  $R(\bar{X}) = L(\bar{X})/OPT(\bar{X})$  measures the penalty paid for serving the tasks in a predetermined order. Here, the service times  $X_i$  are treated as independent identically distributed random variables. Two distributions for  $X_i$ , uniform and exponential, are considered. Bounds on the mean  $ER(\bar{X})$  and on the distribution function  $P[R(\bar{X}) > x]$  are obtained.

**MSC:**

90B35 Deterministic scheduling theory in operations research

Cited in 8 Documents

**Keywords:**

[expected relative performance](#); [list schedule](#); [makespan](#)

**Full Text:** [DOI](#)